

Fully Out-of-Order Choreographies



@dplyukhin

Marco Peressotti



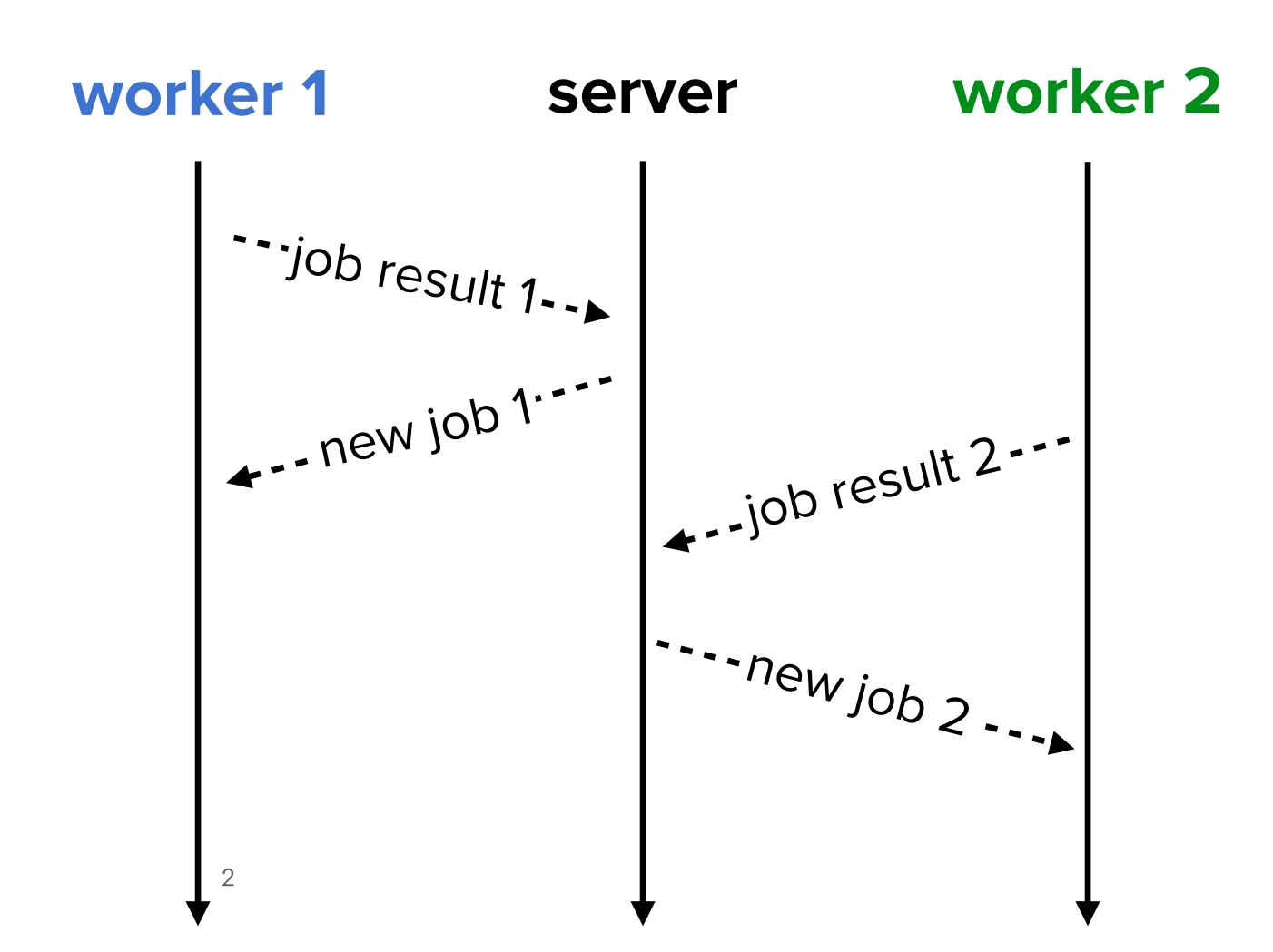
Fabrizio Montesi



what is choreographic programming?



distributed model training



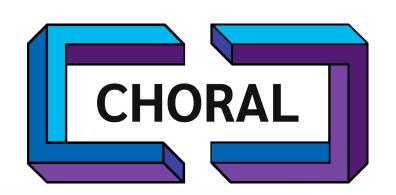
what is choreographic programming?



distributed model training

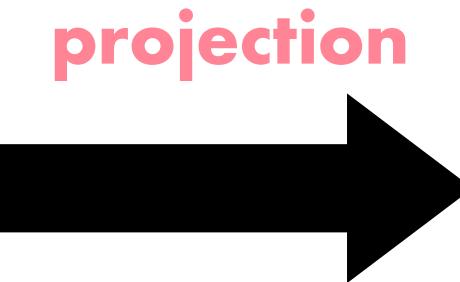
```
// Server.java
                          Result x1 = worker1.recv();
                    184 worker1.send(newJob());
                          Result x2 = worker2.recv();
                          worker2.send(newJob());
     // Worker.java
                                                 // Worker.java
     server.send(results);
                                                 server.send(results);
176
                                            177 Job newJob = server.recv();
     Job newJob = server.recv();
```

choreography







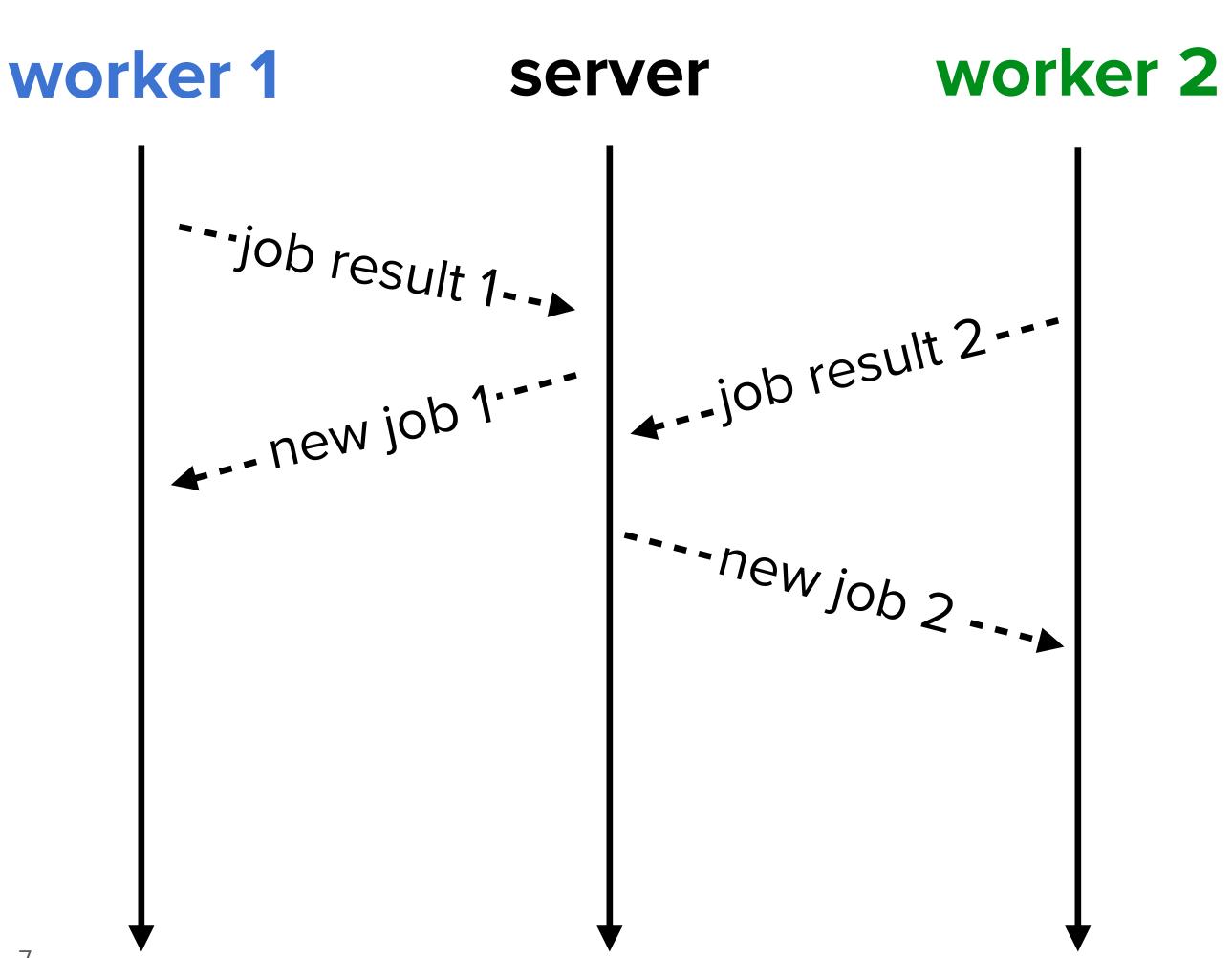


```
endpoint code Java
          // Server.java
          worker
          Result
          worker 175 // Worker.java
                  176 server.send(results);
                 177  Job newJob = server.recv();
  Job newJob = serve
```

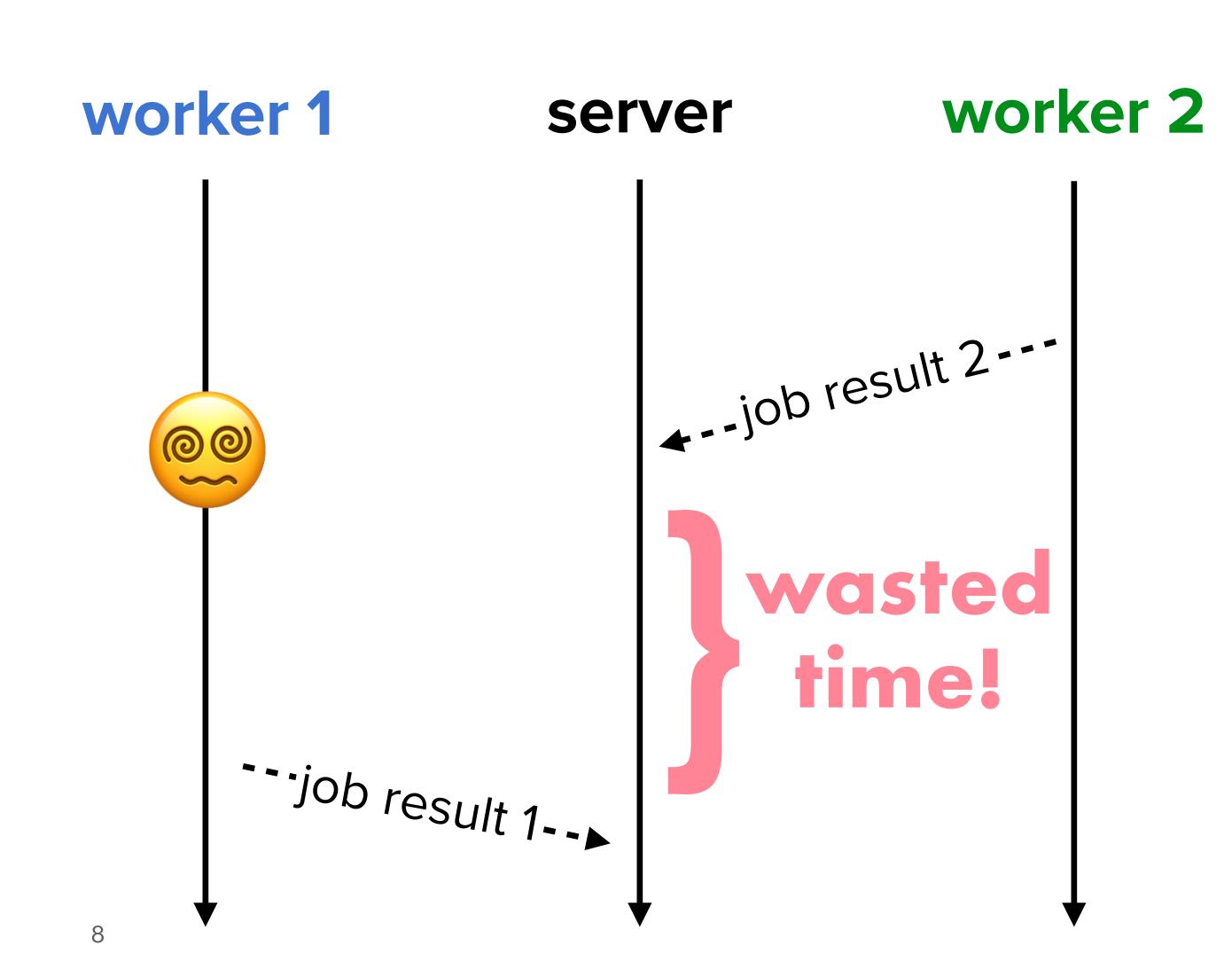
- provably deadlock-free
- no message type errors
- cleaner code

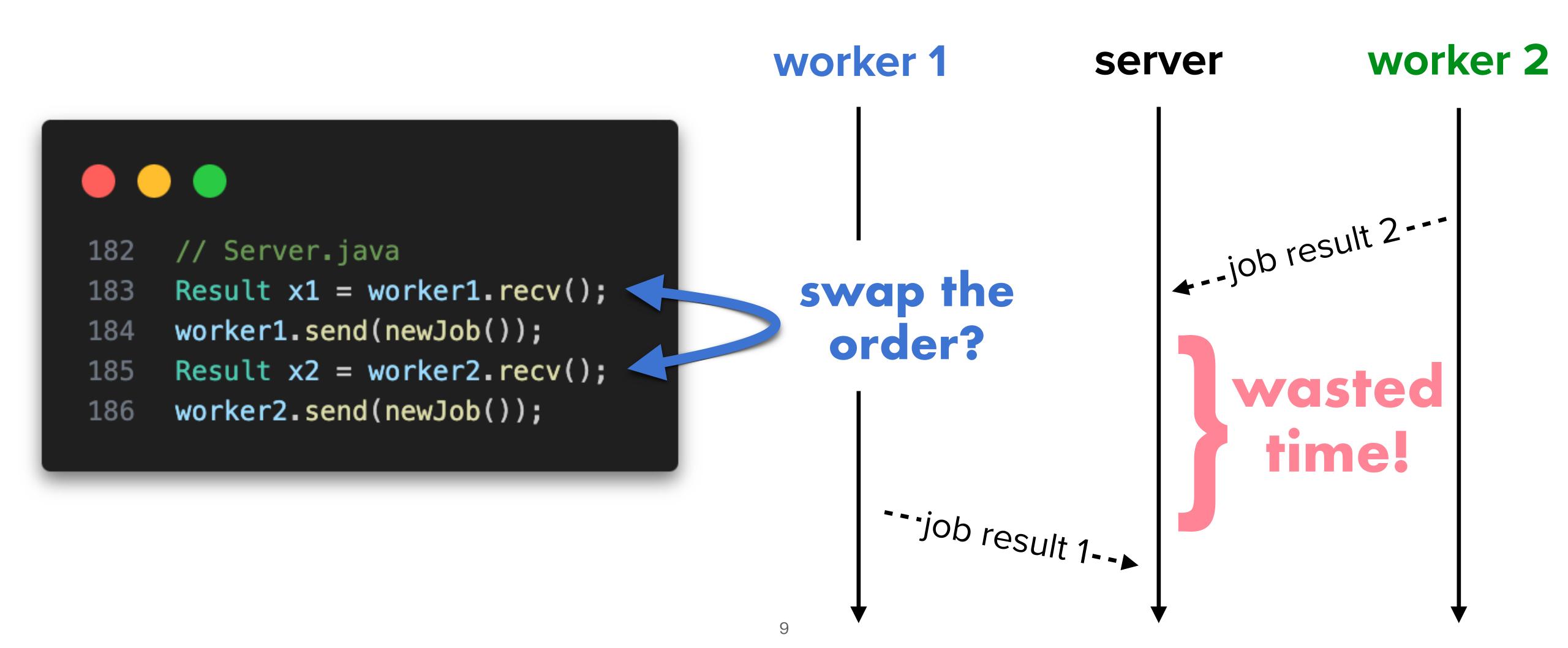
```
182  // Server.java
183  Result x1 = worker1.recv();
184  worker1.send(newJob());
185  Result x2 = worker2.recv();
186  worker2.send(newJob());
```

```
// Server.java
     Result x1 = worker1.recv();
183
     worker1.send(newJob());
184
     Result x2 = worker2.recv();
185
     worker2.send(newJob());
186
```



```
182  // Server.java
183  Result x1 = worker1.recv();
184  worker1.send(newJob());
185  Result x2 = worker2.recv();
186  worker2.send(newJob());
```





```
182  // Server.java
183  Result x2 = worker2.recv();
184  worker2.send(newJob());
185  Result x1 = worker1.recv();
186  worker1.send(newJob());
```

```
worker 1
                                  worker 2
                  server
       ~•job result 1--▶
                     l 4--job result 2---
```

```
182  // Server.java
183  Future<Result> f1 = worker1.recv();
184  f1.andThen(x -> worker1.send(newJob()));
185  Future<Result> f2 = worker2.recv();
186  f2.andThen(x -> worker2.send(newJob()));
```

we need

out-of-order

processes!

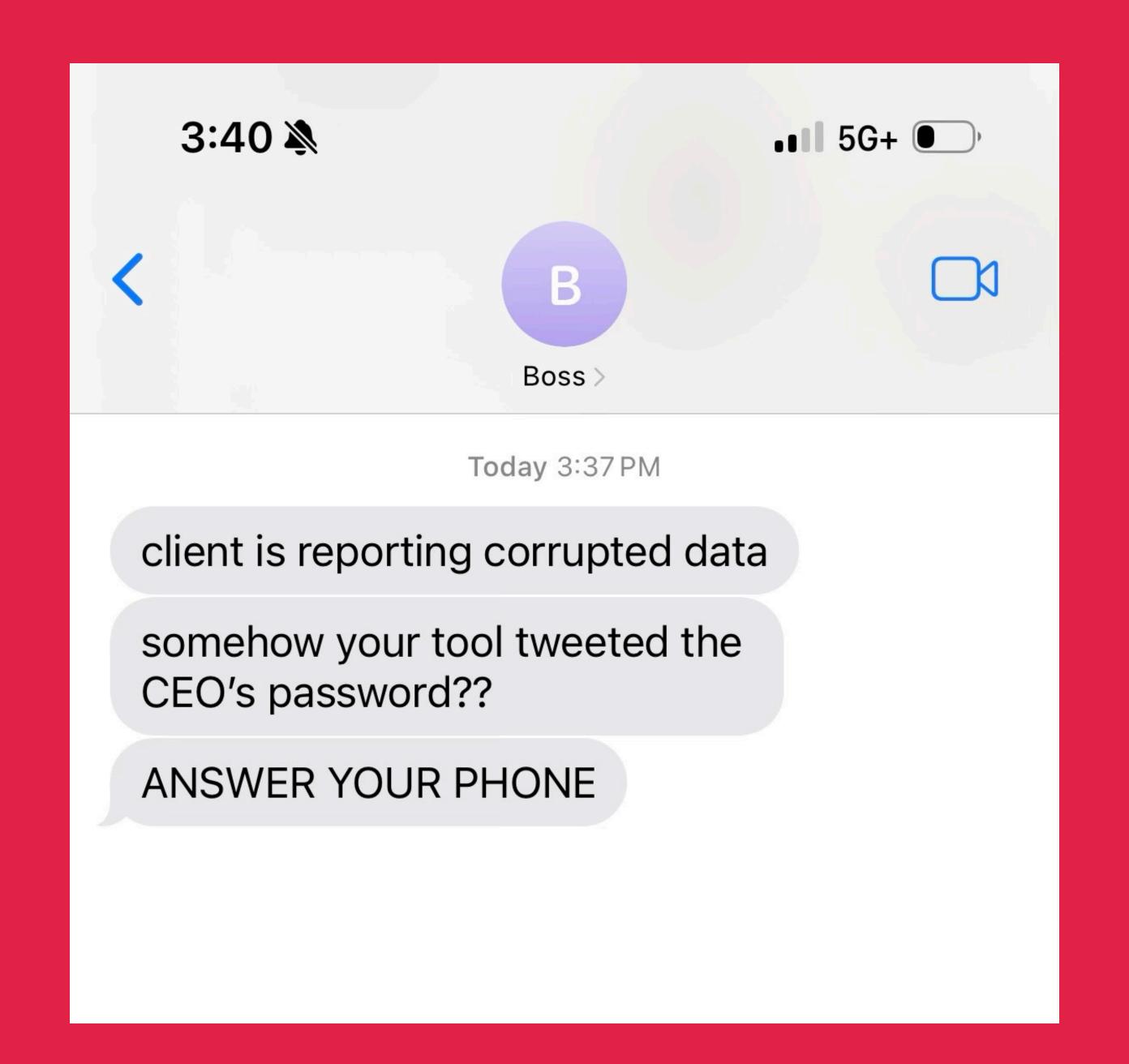
easy!

- 1. patch the compiler
- 2. deploy to production
 - 3. happy friday



CODE RED





this talk

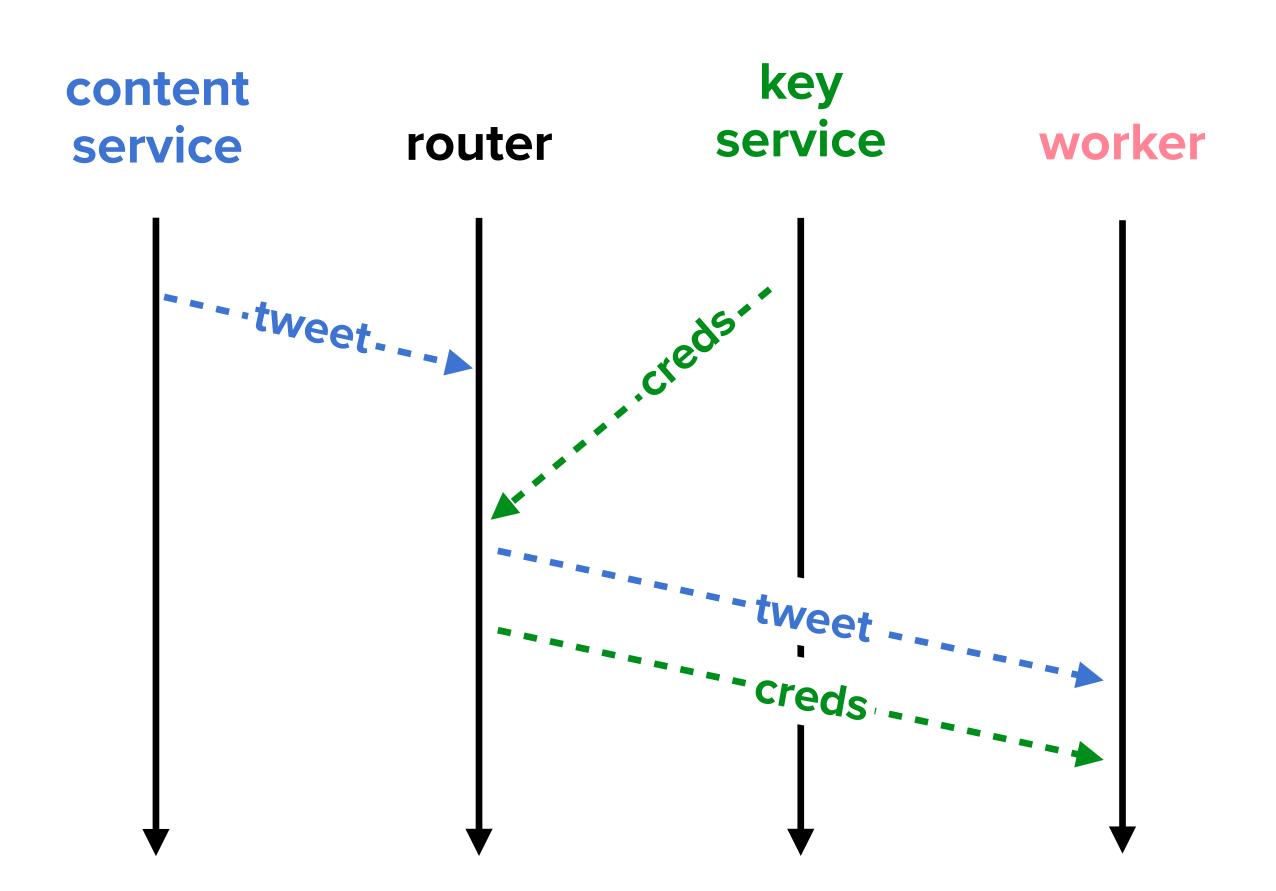




what went wrong?

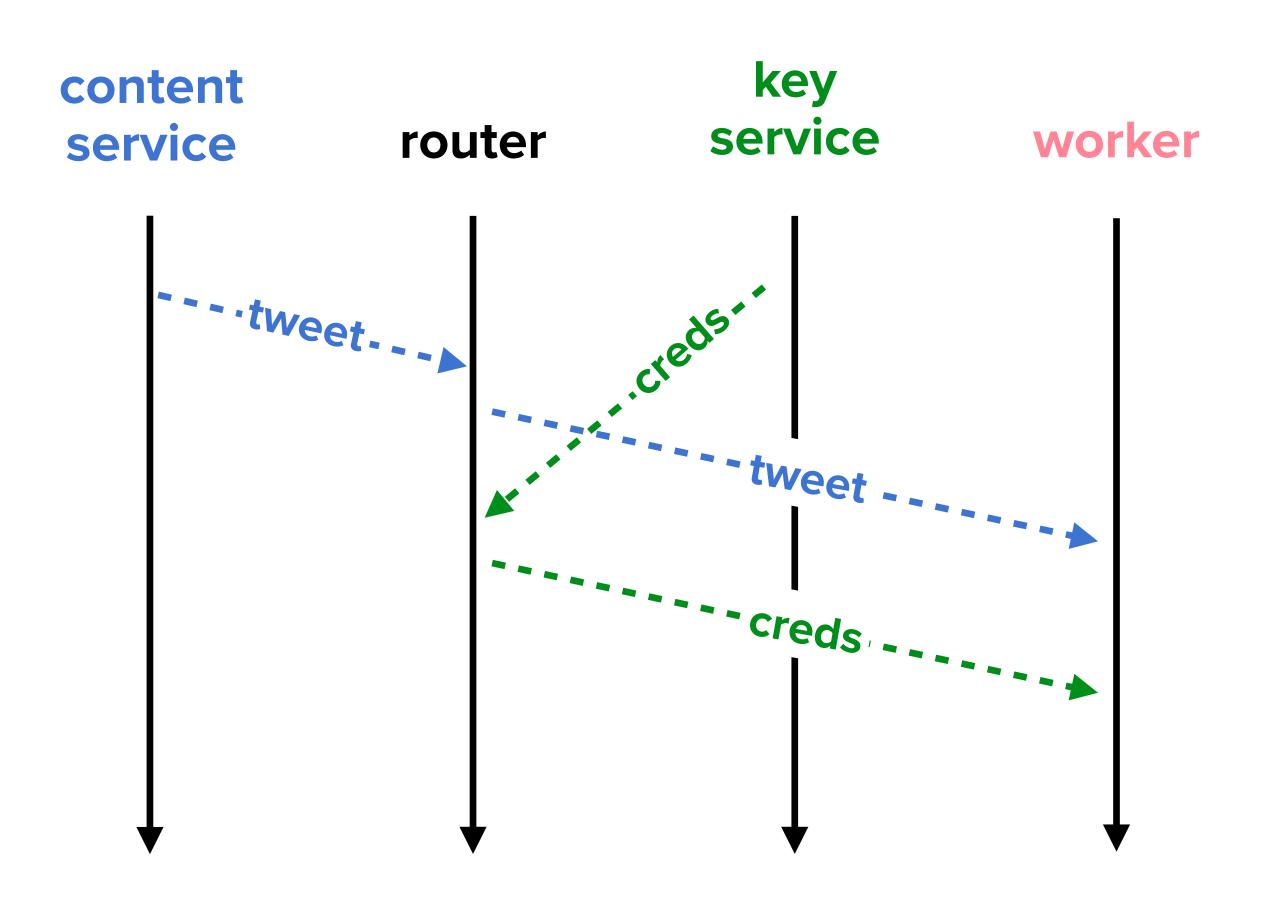
1 // TweetChoreography.ch
2 cs.tweet -> router.tweet;
3 ks.creds -> router.creds;
4 router.tweet -> worker.tweet;
5 router.creds -> worker.creds;
6 worker.post(tweet, creds);

choreography



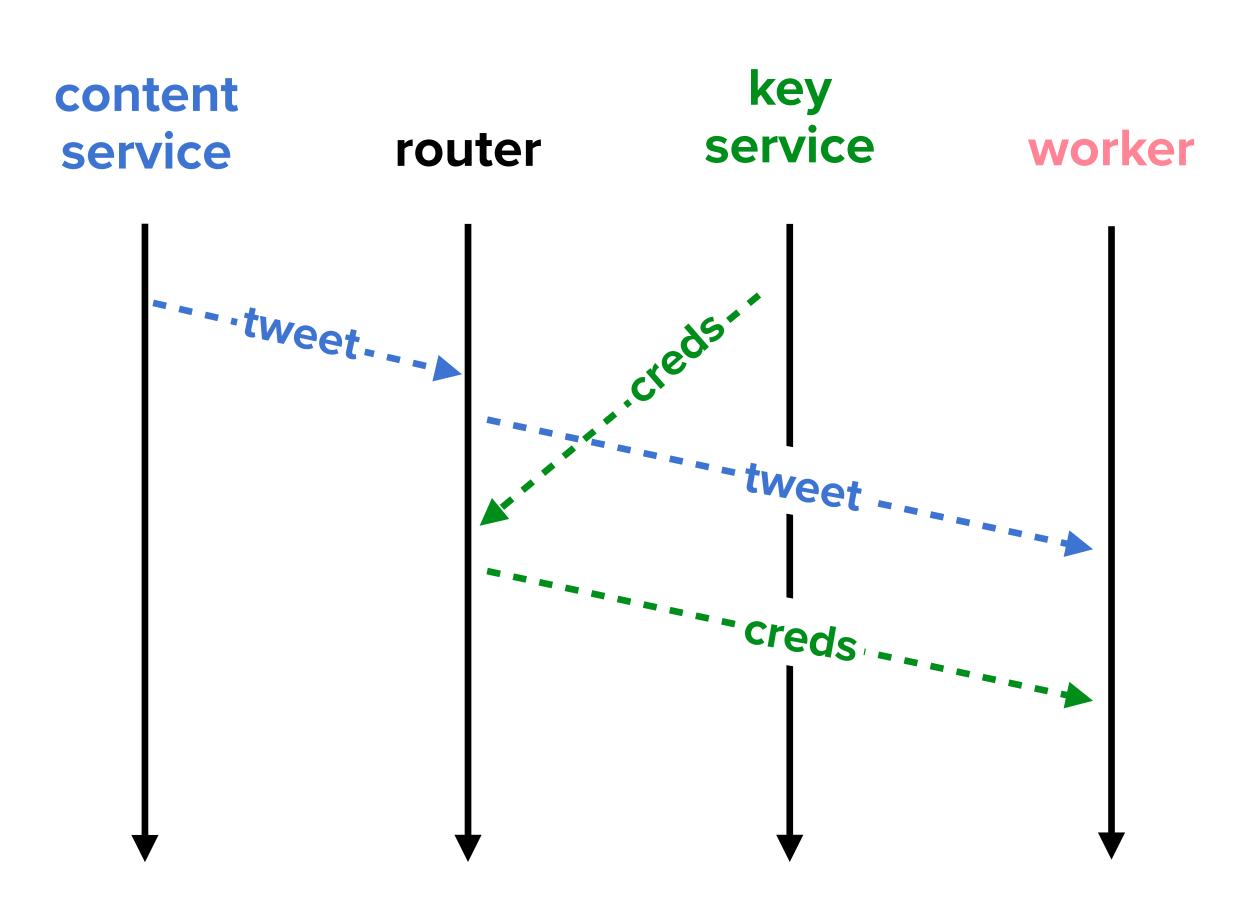
1 // TweetChoreography.ch
2 cs.tweet -> router.tweet;
3 ks.creds -> router.creds;
4 router.tweet -> worker.tweet;
5 router.creds -> worker.creds;
6 worker.post(tweet, creds);

choreography



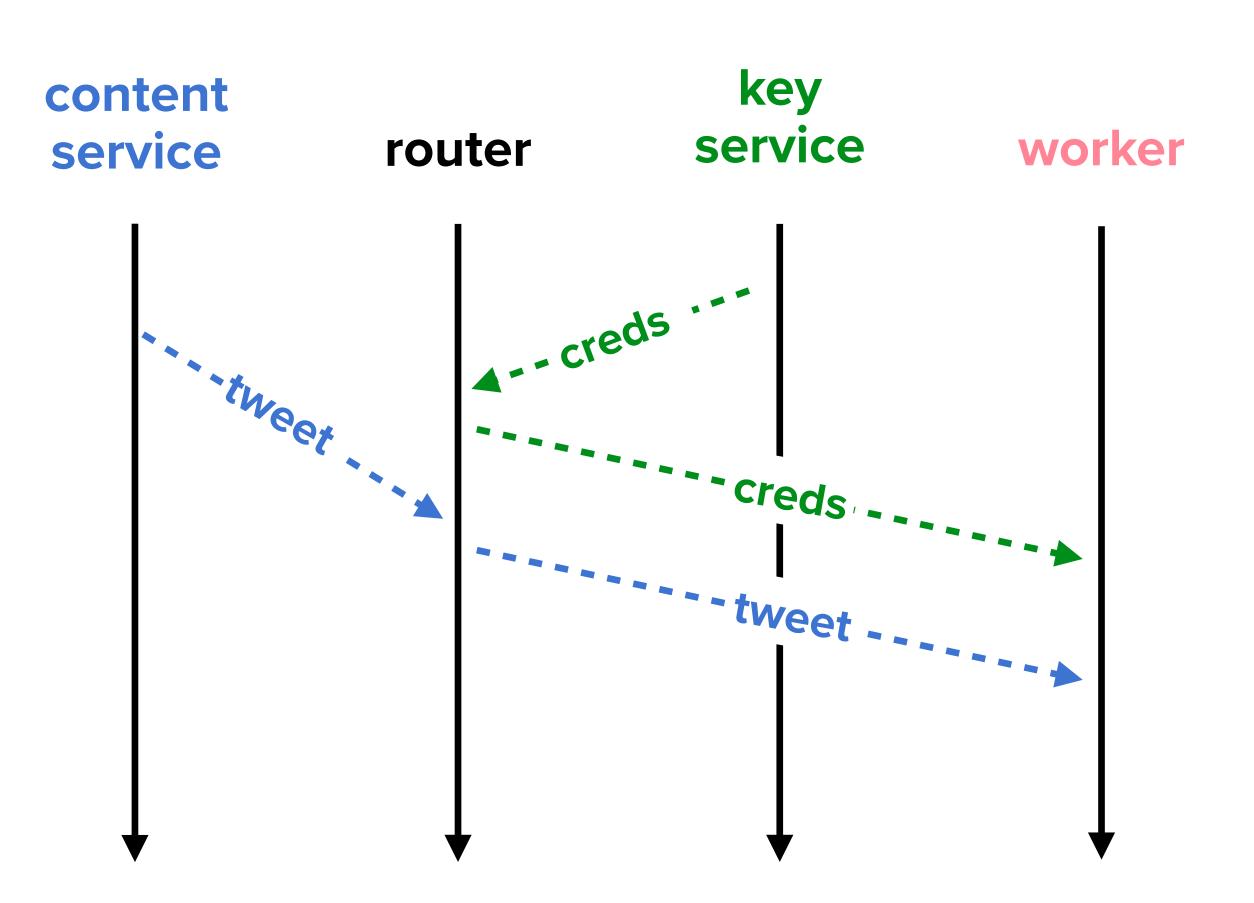
1 // Router.java
2 Future<String> tweet = cs.recv();
3 Future<String> creds = ks.recv();
4 tweet.andThen(t -> worker.send(t));
5 creds.andThen(c -> worker.send(c));

compiled endpoint code



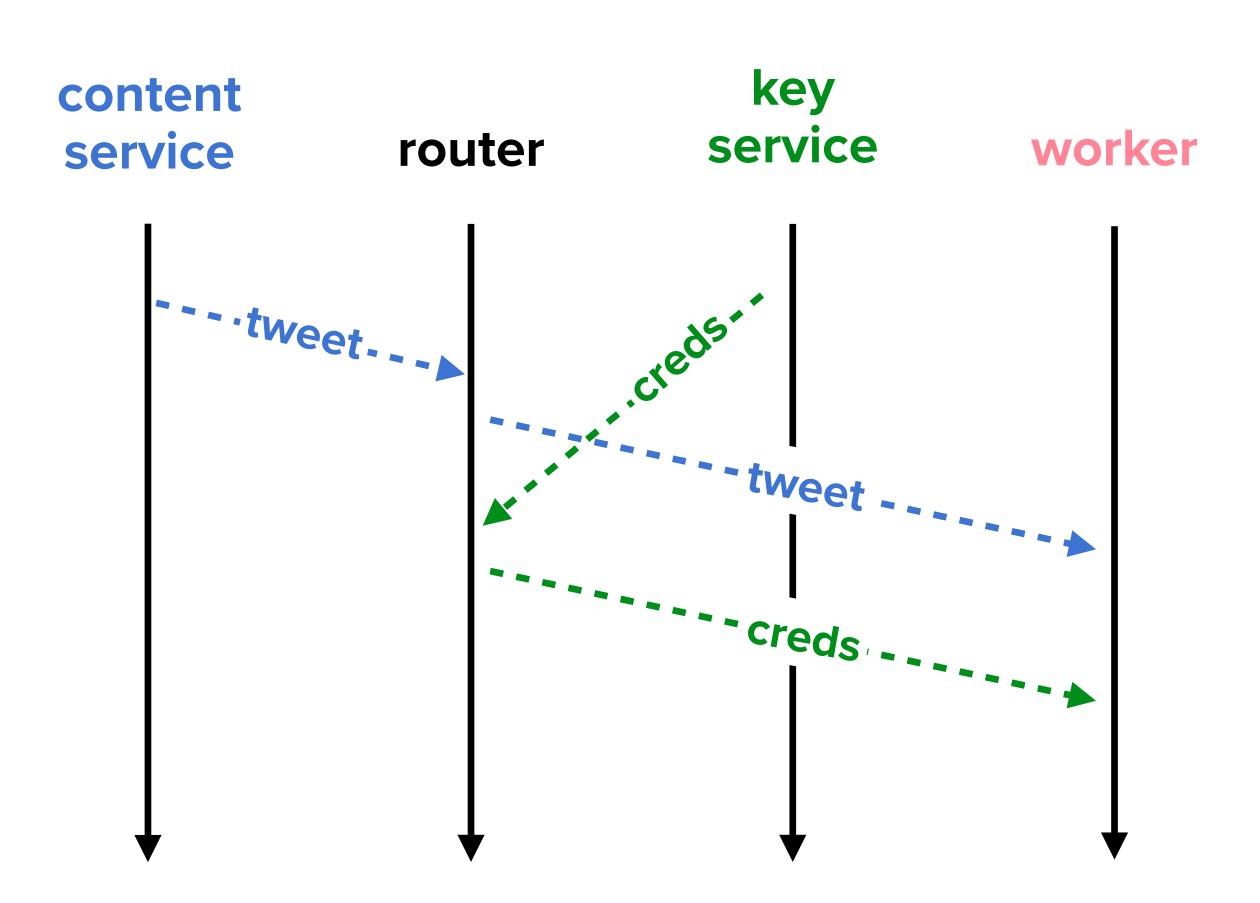
1 // Router.java
2 Future<String> tweet = cs.recv();
3 Future<String> creds = ks.recv();
4 tweet.andThen(t -> worker.send(t));
5 creds.andThen(c -> worker.send(c));

compiled endpoint code



1 // Router.java
2 Future<String> tweet = cs.recv();
3 Future<String> creds = ks.recv();
4 tweet.andThen(t -> worker.send(t));
5 creds.andThen(c -> worker.send(c));

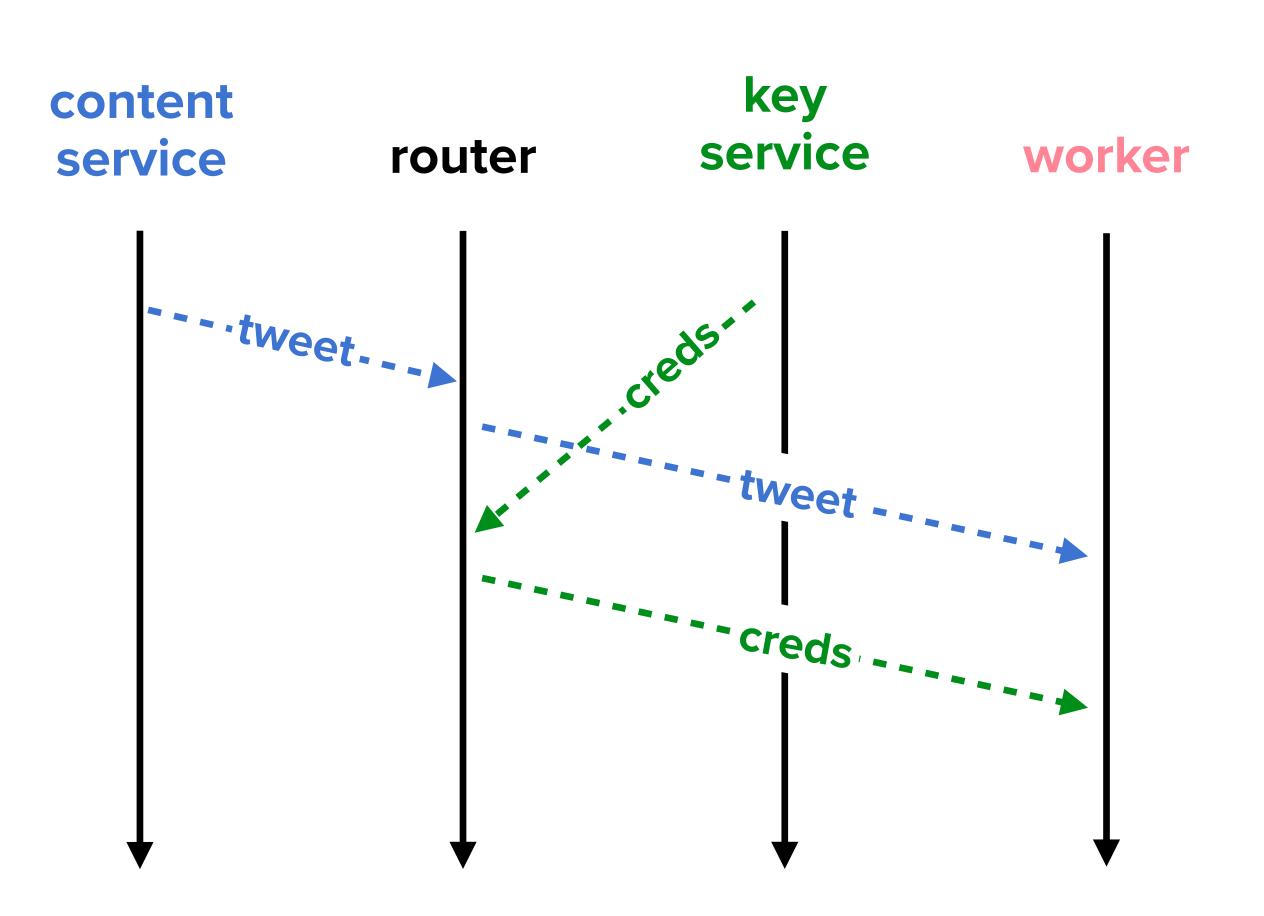
compiled endpoint code



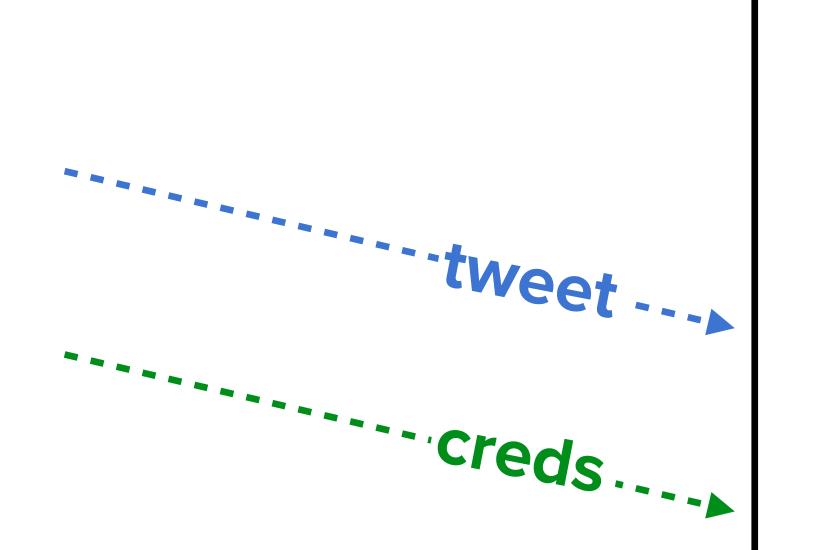
the **content service** generates tweets the **key service** manages login info the **worker** calls the Twitter API

```
// Router.java
Future<String> tweet = cs.recv();
Future<String> creds = ks.recv();
tweet.andThen( t -> worker.send(t) );
creds.andThen( c -> worker.send(c) );
// Worker.java
Future<String> tweet = router.recv();
Future<String> creds = router.recv();
post(tweet.get(), creds.get());
```

compiled endpoint code



worker



```
1  // Worker.java
2  Future<String> tweet = router.recv();
3  Future<String> creds = router.recv();
4  post(tweet.get(), creds.get());
```

communication integrity worker violation (CIV) ·-··creds. // Worker.java Future<String> tweet = router.recv(); Future<String> creds = router.recv(); post(tweet.get() creds.get());

router

worker

communication integrity violation (CIV)

```
(2, creds) --- I
    --(1, tweet) --->
```

```
1  // Worker.java
2  Future<String> tweet = router.re(v(1);
3  Future<String> creds = router.re(v(2);
4  post(tweet.get(), creds.get());
```

router

worker

communication integrity violation (CIV)

```
`'(2, creds) ---
```

```
1  // Worker.java
2  Future<String> tweet = router.re(v(1);
3  Future<String> creds = router.re(v(2);
4  post(tweet.get(), creds.get());
```

router

worker

communication integrity violation (CIV)

```
(2, creds) --- I
    --(1, tweet) --->
```

```
1  // Worker.java
2  Future<String> tweet = router.re(v(1);
3  Future<String> creds = router.re(v(2);
4  post(tweet.get(), creds.get());
```

but wait, there's more!

big takeaway #1

You can prevent CIVs inside a choreography with statically unique IDs!

but wait, there's more!

online shopping checkout



billing cart 1 exchange cart 2 dan charge(dan, \$1000) marco 32

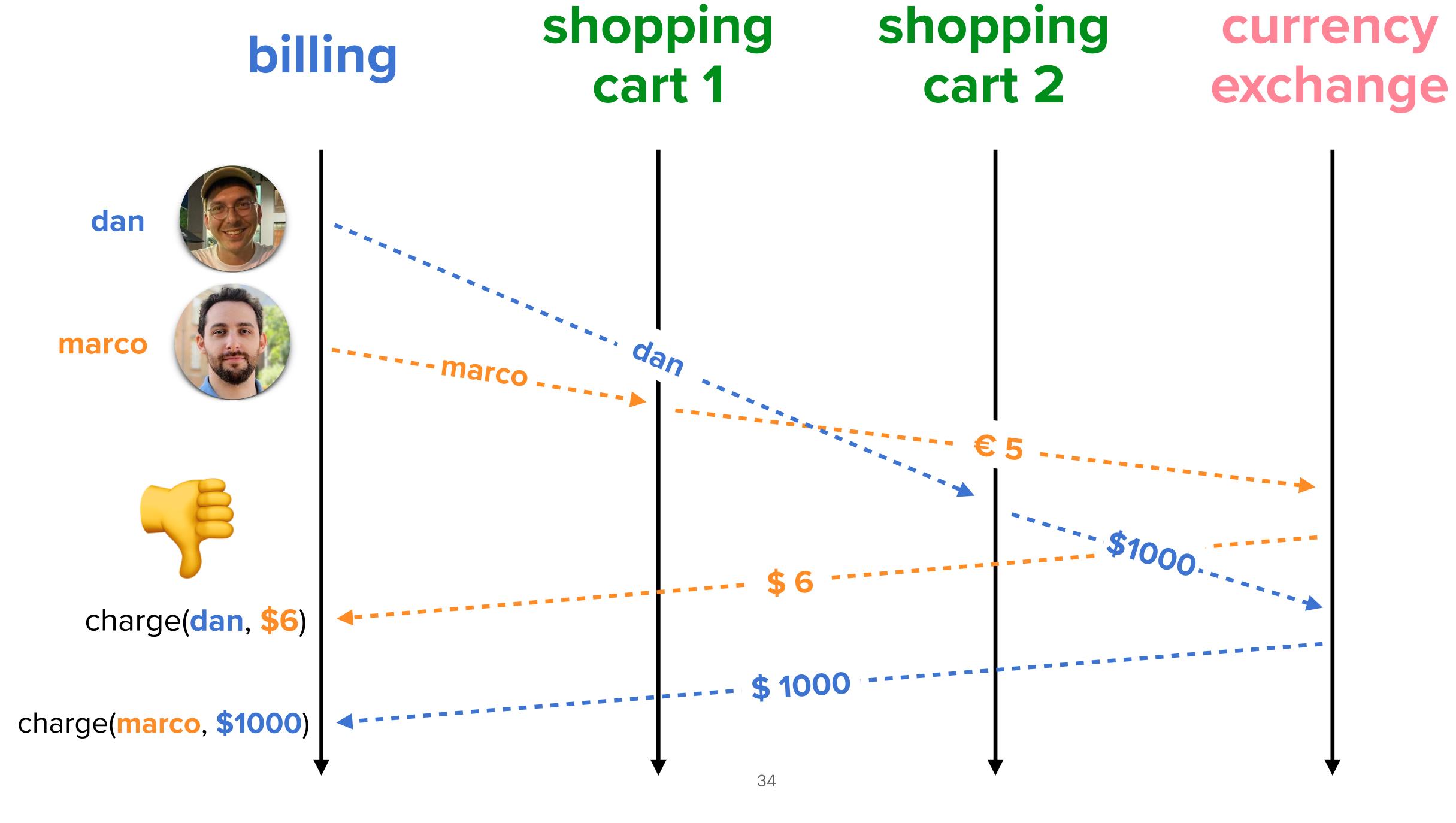
shopping

shopping

currency

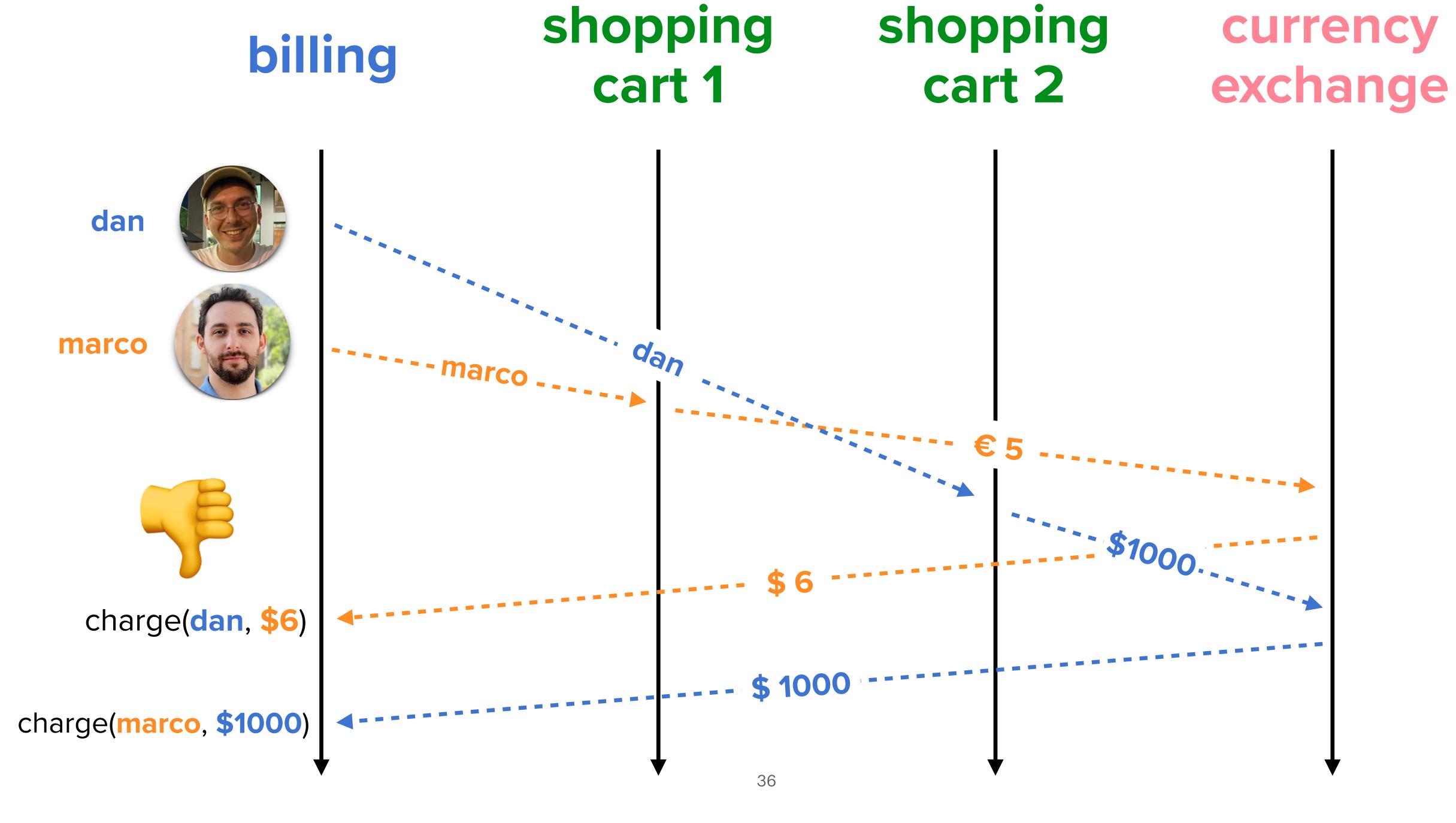
shopping shopping currency billing cart 1 cart 2 exchange dan marco charge(dan, \$1000) charge(marco, \$6)

33



shopping shopping currency billing cart 1 cart 2 exchange dan marco charge(dan, \$1000) charge(marco, \$6)

35



billing exchange cart 1 cart 2 dan marco charge(marco, \$6) charge(dan, \$1000) 37

shopping

shopping

currency

big takeaway #2

You can prevent CIVs between choreographies with dynamically unique session tokens!

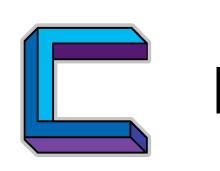
the paper



formal model: session tokens without synchronization

proofs: deadlock-freedom, bisimulation, communication integrity

performance: microbenchmarks, model serving









conclusion



big takeaway #1

read the paper!

You can prevent CIVs inside a choreography with statically unique IDs!

big takeaway #2

You can prevent CIVs between choreographies with dynamically unique session tokens!